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STOCHASTIC NEURAL NETWORKS

by

Eugene Wong

Memorandum No. UCB/ERL M89/9

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Stochastic Neural Networks

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1. Hopfield Networks

It is well known [1, 2] that a neural network can be used to compute the minimum of a function $E(x)$ defined on a hypercube $[0, 1]^n$ as follows. Let $v_i(t)$ be the state at node i at time t and set

$$v_i(t) = g(u_i(t)) \quad (1.1)$$

$$\frac{d u_i(t)}{dt} = -E_i(v(t)) \quad (1.2)$$

where $E_i(v) = \frac{\partial}{\partial v_i} E(v)$ and g is an increasing function. The special case of a quadratic E

$$E(v) = -\frac{1}{2} \sum_{i,j} W_{ij} v_i v_j - \sum_i \theta_i v_i \quad (1.3)$$

where we assume $W_{ij} = W_{ji}$, results in

$$E_i(v) = -\sum_j W_{ij} v_j - \theta_i \quad (1.4)$$

which is particularly well suited for realization as an analog integrated circuit.

As Hopfield and Tank [3] and others [e.g., [4]] have shown, a variety of computational problems of considerable complexity can be reduced to computing the global minimum of a quadratic function. With current technology, a network with several hundred nodes and with the dynamics given by equation (1.1) through (1.3) can probably be built on a single chip. The potential for simple and fast computation thus created is exciting indeed.

However, (1.1) and (1.2) represent essentially a gradient-descent method for minimization, and such methods do not always reach a global minimum. To see this we write using (1.1) and (1.2)

$$\begin{aligned} \frac{d}{dt} E(v(t)) &= \sum_i E_i(v(t)) \frac{d v_i(t)}{dt} \\ &= \sum_i E_i(v(t)) g'(u_i(t)) \frac{d u_i(t)}{dt} \end{aligned} \quad (1.5)$$

$$= - \sum_i g'(u_i(t)) E_i^2(v(t)) \leq 0$$

which shows that E is decreasing but not *strictly* decreasing and the equilibrium reached may only be a local minimum.

2. Boltzmann Machines

For some problems it is sufficient to restrict v_i to binary values, say $v_i = 0, 1$. The collection of the states at all the nodes of the network $v(t) = \{v_i(t)\}$ now takes values in $\{0,1\}^n$, and $v(t)$ will be called the *configuration* of the network at time t . A Boltzmann machine [5] is a network where $v(t)$ is a $\{0, 1\}^n$ valued discrete-time Markov chain with state transitions defined as follows:

For each $v \in \{0, 1\}^n$ define a neighborhood $N(v) \in \{0, 1\}^n$. We assume that $v' \in N(v) \Rightarrow v \in N(v')$ and $v \notin N(v)$. At time $t + 1$ we choose a $v' \in N(v(t))$ at random (say with equal probabilities) and set

$$\begin{aligned} v(t+1) &= v' && \text{with probability } p(\Delta E) \\ &= v(t) && \text{with probability } 1 - p(\Delta E) \end{aligned} \quad (2.1)$$

where ΔE is given by

$$\Delta E = E(v') - E(v(t)) \quad (2.2)$$

and $p(\Delta E)$ is a probability of the form

$$p(\Delta E) = e^{-\frac{1}{2T} \Delta E} f(|\Delta E|) \quad (2.3)$$

for some decreasing function f . Familiar examples include

$$\begin{aligned} p(\Delta E) &= \min(1, e^{-\frac{1}{T} \Delta E}) \\ &= e^{-\frac{1}{2T} \Delta E} e^{-\frac{1}{2T} |\Delta E|} \end{aligned}$$

and

$$p(\Delta E) = \frac{1}{1 + e^{\frac{1}{T} \Delta E}} = \frac{2e^{-\frac{1}{2T} \Delta E}}{\cosh(\frac{\Delta E}{2T})}$$

It is clear that $v(\cdot)$ has a bias for moving in the direction of decreasing ΔE , but will move with nonzero probabilities even for increasing ΔE .

The one-step transition probability is given by

$$\begin{aligned} P(v|v_0) &= \text{Prob}(v(t+1) = v | v(t) = v_0) \\ &= \frac{1}{|N(v_0)|} p(E(v) - E(v_0)) \quad v \in N(v_0) \\ &= 0 \quad v \neq v_0, v \notin N(v_0) \end{aligned} \tag{2.4}$$

and

$$P(v_0|v_0) = 1 - \frac{1}{|N(v_0)|} \sum_{v' \in N(v_0)} p(E(v') - E(v_0))$$

where $|N(v_0)|$ denotes the cardinality of $N(v_0)$. With p given by (2.3), we have

$$\begin{aligned} \sum_{v_0} P(v|v_0) e^{-\frac{1}{T} E(v_0)} |N(v_0)| \\ &= P(v|v) e^{-\frac{1}{T} E(v)} |N(v)| + \sum_{v_0 \in N(v)} P(v|v_0) e^{-\frac{1}{T} E(v_0)} |N(v_0)| \\ &= |N(v)| e^{-\frac{1}{T} E(v)} - \sum_{v' \in N(v)} f(|E(v') - E(v)|) e^{-\frac{1}{2T} [E(v) + E(v')]} \\ &\quad + \sum_{v_0 \in N(v)} f(|E(v) - E(v_0)|) e^{-\frac{1}{2T} [E(v) + E(v_0)]} \\ &= |N(v)| e^{-\frac{1}{T} E(v)} \end{aligned}$$

It follows that a probability distribution of the form

$$P(v(t) = v) = K |N(v)| e^{-\frac{1}{T} E(v)}$$

is left invariant by the transitions.

If, in addition, the Markov chain is irreducible, i.e., every v can be reached from any initial configuration v_0 , then

$$P(v(t) = v | v(t_0) = v_0) \xrightarrow{(t-t_0) \rightarrow \infty} K |N(v)| e^{-\frac{1}{T} E(v)} \quad (2.5)$$

where K is the normalizing constant. If we assume $|N(v)|$ is independent of v , then the stationary distribution is simply

$$P(v) = \frac{1}{Z} e^{-\frac{1}{T} E(v)} \quad (2.6)$$

where $Z = \sum_v e^{-\frac{1}{T} E(v)}$ is called the partition function in statistical mechanics.

Equation (2.6) is called the Gibbs or Boltzmann distribution, and the Markov chain $v(t)$ a *Gibbs field*. A network with such a $v(t)$ has been called a *Boltzmann machine* [5]. It could also have been called a Gibbs machine.

3. Simulated Annealing

Because the stationary distribution of a Boltzmann machine is given by

$$P(v) = \frac{1}{Z} e^{-\frac{1}{T} E(v)} \quad (3.1)$$

the peaks of $P(v)$ coincide with the minima of $E(v)$. As the parameter T (temperature) decreases to zero, $P(v)$ will approach a set of Dirac δ -functions at the global minima of $E(v)$. This is the principle on which *simulated annealing* is based [6].

Suppose that we choose a sequence $\{T_k\}$ decreasing to 0 sufficiently slowly so that for large k , $v(k)$ is distributed approximately according to

$$P_k(v) = \frac{1}{Z_k} e^{-\frac{1}{T_k} E(v)}$$

Then we would expect $v(k)$ to converge to a global minimum. This is indeed the case for T_k of the form

$$T_k = \frac{c}{\ln(1+k)} \quad (3.2)$$

where c is a ‘‘sufficiently large’’ constant [7].

Simulated annealing can be extended to the continuous variable case. This is done with the Langevin algorithm [8, 9], which is defined by a set of stochastic differential equations of the form

$$dv_i(t) = -E_i(v(t)) dt + \sqrt{2T} dW_i(t) \quad (3.3)$$

where $E_i = \frac{\partial}{\partial v_i} E$ as before and $\{W_i\}$ is a set of independent Wiener processes. The goal, once again, is to get a stationary distribution for $v(t)$ characterized by the density function

$$p(v) = \frac{1}{Z} e^{-\frac{1}{T} E(v)} \quad (3.4)$$

However, for $v \in \mathbb{R}^n$ there may be no density function of this form since $\exp(-\frac{1}{T} E)$ may not be integrable. For $v \in [0, 1]^n$, (3.3) also does not guarantee a stationary density of the form (3.4) without imposing a ‘‘reflecting boundary’’ condition at every boundary $v_i = 0, 1$ [8].

4. Diffusion Machines: Stochastic Hopfield Networks

We propose a scheme that is a modification of both the Langevin algorithm and the Hopfield network. Suppose that we inject noise in a Hopfield network so that at the i th node the equations of dynamics are now given by (c.f. (1.1) and (1.2)):

$$v_i(t) = g(u_i(t)) \quad (4.1)$$

$$du_i(t) = -E_i(v(t)) dt + \alpha_i(u(t)) dW_i(t) \quad (4.2)$$

where (4.2) is a stochastic differential equation of the Ito type [10]. As in the Langevin algorithm, $\{W_i\}$ are independent Wiener processes.

The question we now pose is the following: Can α_i 's be found so that $v(t)$ is a stationary Markov process with the following stationary density?

$$p_0(v) = \frac{1}{Z} e^{-\frac{1}{T} E(v)} \quad (4.3)$$

The answer is surprisingly simple, and unique. The required α_i is given by

$$\alpha_i(u(t)) = \sqrt{\frac{2T}{g'(u_i(t))}} \quad (4.4)$$

so that (4.2) becomes

$$d u_i(t) = -E_i(v(t)) dt + \sqrt{\frac{2T}{g'(u_i(t))}} d W_i(t) \quad (4.5)$$

Furthermore, if we denote

$$f(x) = g'(g^{-1}(x)) \quad (4.6)$$

then $v_i(t)$ satisfies the stochastic differential equation.

$$d v_i(t) = -f(v_i(t)) E_i(v(t)) dt + T f'(v_i(t)) dt + \sqrt{2T f(v_i(t))} d W_i(t) \quad (4.7)$$

If $f(x) = 0$ at $x = 0, 1$, then stationarity of $v(t)$ is assured. If not, a reflecting boundary is needed at each $v_i = 0, 1$.

To derive (4.4), we first note that with a smooth g , we can use the Ito differentiation formula [11] and derive a set of stochastic differential equations for v_i which are of the form

$$d v_i(t) = m_i(v(t)) dt + \sigma_i(v(t)) d W_i(t) \quad (4.8)$$

The transition density $p(v, t | v_0, t_0)$ of $v(t)$ must satisfy the Fokker-Planck equation

$$\frac{\partial p}{\partial t} = \sum_i \frac{\partial}{\partial v_i} \left[\frac{1}{2} \frac{\partial}{\partial v_i} (\sigma_i^2 p) - m_i p \right] \quad (4.9)$$

It follows that if (4.3) is to be the stationary density then we must have

$$\sum_i \frac{\partial}{\partial v_i} \left[\frac{1}{2} \frac{\partial}{\partial v_i} (\sigma_i^2 p_0) - m_i p_0 \right] = 0 \quad (4.10)$$

which is satisfied if m_i and σ_i^2 satisfy

$$m_i = \frac{1}{2} \sigma_i^2 \left(-\frac{1}{T} E_i(v) \right) + \frac{\partial}{\partial v_i} \left(\frac{1}{2} \sigma_i^2 \right) \quad (4.11)$$

Since (4.8) is derived from (4.1) and (4.2) using the Ito differentiation formula, we have

$$d v_i(t) = g'(u_i(t)) d u_i(t) + \frac{1}{2} g''(u_i(t)) \alpha_i^2(u_i(t)) dt \quad (4.12)$$

Comparing (4.12) and (4.8), we get

$$\sigma_i(g(u)) = g'(u_i) \alpha_i(u) \quad (4.13)$$

and

$$m_i(v) = -f(v_i) E_i(v) + \frac{1}{2} (\ln f)'(v_i) \sigma_i^2(v) \quad (4.14)$$

where f is given by

$$f(x) = g'(g^{-1}(x)) \quad (4.6)$$

Comparing (4.6) with (4.11) now yields

$$\frac{1}{2} \sigma_i^2(v) = T f(v_i) \quad (4.15)$$

which is a relationship of great simplicity.

A network with dynamics governed by (4.1) and (4.5) (equivalently (4.7)) will be called a *diffusion machine*. We propose that it be used as the basis for studying simulated annealing and machine learning [5]. As a neural computing system, it has a number of important advantages. First, it is quite general. There is no need to assume that the minimum occurs at a corner of the cube. Second, it allows the non-linearity g to play a stabilizing role in ensuring stationarity. Finally and most importantly, it is well suited for direct circuit implementation, thus providing potentially much faster computation for both annealing and machine learning.

5. An Example

A favorite choice of g is

$$g(x) = \frac{1}{2} (1 + \tanh \frac{x}{a}) \quad (5.1)$$

which yields

$$g'(x) = \frac{1}{2a} (1 - \tanh^2 \frac{x}{a}) = \frac{1}{2a} (\cosh \frac{x}{a})^{-2} \quad (5.2)$$

and

$$\begin{aligned} f(y) &= g'(g^{-1}(y)) = \frac{1}{2a} [1 - (2y - 1)^2] \\ &= \frac{2}{a} y(1 - y) \end{aligned} \quad (5.3)$$

Equations (4.5) and (4.7) now become

$$d u_i(t) = -E_i(v(t)) dt + 2\sqrt{aT} \cosh \left[\frac{u_i(t)}{a} \right] d W_i(t) \quad (5.4)$$

$$\begin{aligned} d v_i(t) &= -\frac{2}{a} v_i(t) [1 - v_i(t)] E_i(v(t)) dt + \frac{2T}{a} [1 - 2v_i(t)] dt \\ &\quad + 2\sqrt{\frac{T}{a}} v_i(t) [1 - v_i(t)] d W_i(t) \end{aligned} \quad (5.5)$$

As a second example, consider the case

$$\begin{aligned} g(x) &= x \quad , \quad 0 \leq x \leq 1 \\ &= 1 \quad , \quad x > 1 \\ &= 0 \quad , \quad x < 0 \end{aligned}$$

This example corresponds to the Langevin algorithm as considered in [8]. Because $g^{-1}(v)$ does not exist in this case, this example is not really a diffusion machine. If it is to be considered at all, reflecting boundaries at $v_i = 0, 1$, are required [8].

6. Rate of Convergence

A diffusion machine can be used with a cooling schedule $\{T_k\}$ to achieve simulated annealing. For that purpose an estimate of the rate at which

$$p(v, t | v_0, t_0) \rightarrow p_0(v)$$

is needed. Diffusion theory provides a powerful approach to such estimates (c.f., [9]).

Under the assumption $f(x) = 0$, $x = 0, 1$, we can write [??]

$$p(v, t | v_0, 0) = p_0(v) \sum_{\lambda} e^{-\lambda t} \psi_{\lambda}(v) \psi_{\lambda}(v_0) \quad (6.1)$$

where λ are the eigenvalues, and ψ_{λ} the normalized eigen functions, of the equation

$$T \sum_i \frac{\partial}{\partial v_i} \left[p_0(v) f(v_i) + \frac{\partial \psi_{\lambda}(v)}{\partial v_i} \right] + \lambda p_0(v) \psi_{\lambda}(v) = 0 \quad (6.2)$$

It can be shown that $\lambda = 0$ is the smallest eigenvalue with $\psi_0(v) = 1$, and the eigenvalues can be ordered

$$0 = \lambda_0 < \lambda_1 \leq \lambda_2 \dots$$

Using (6.2) for ψ_1 , multiplying each term by ψ_1 , and integrating, we get

$$\lambda_1 \int_{[0,1]^p} p_0(v) \psi_1^2(v) dv = T \sum_i \int_{[0,1]^p} p_0(v) f(v_i) \left[\frac{\partial \psi_1(v)}{\partial v_i} \right]^2 dv$$

Since ψ_1 is normalized and orthogonal to $\psi_0 = 1$, we get

$$\lambda_1 = \min_{\psi} T \int_{[0,1]^p} p_0(v) \sum_i f(v_i) \left[\frac{\partial \psi(v)}{\partial v_i} \right]^2 dv \quad (6.3)$$

subject to the conditions

$$\int_{[0,1]^p} p_0(v) \psi^2(v) dv = 1 \quad (6.4)$$

and

$$\int_{[0,1]^p} p_0(v) \psi(v) dv = 0 \quad (6.5)$$

It is clear that $\lambda_1 > 0$, and its dependence on T and g can be studied via (6.3).

From (6.1) we get

$$\begin{aligned} |p(v, t | v_0, 0) - p_0(v)| &= \left| \sum_{k=1}^{\infty} e^{-\lambda_k t} \psi_k(v) \psi_k(v_0) \right| \\ &= e^{-\lambda_1 t} \left| \sum_{k=1}^{\infty} e^{-(\lambda_k - \lambda_1)t} \psi_k(v) \psi_k(v_0) \right| \end{aligned}$$

If we denote

$$K(v) = e^{\lambda_1} |p(v, 1 | v, 0) - p_0(v)|$$

then for $t > 1$

$$\begin{aligned} |p(v, t | v_0, 0) - p_0(v)| &\leq e^{-\lambda_1 t} \sqrt{k(v)k(v_0)} \\ &\leq e^{-\lambda_1 t} \sup_v k(v) \end{aligned} \quad (6.6)$$

which gives an estimate of the rate of convergence of the transition density to the equilibrium distribution.

7. Analog Realization

Equations (4.5) and (4.7) are Ito equations. To realize them in analog circuits using Gaussian wide-band noise requires a correction term [10]. With the correction term, (4.5) and (4.7) can be rewritten as

$$\dot{u}_i(t) = -E_i(v(t)) + \frac{T}{2} \frac{g''(u_i(t))}{[g'(u_i(t))]^2} + \sqrt{\frac{2T}{g'(u_i(t))}} \eta_i(t) \quad (7.1)$$

$$\dot{v}_i(t) = -f(v_i(t)) E_i(v(t)) + \frac{T}{2} f'(v_i(t)) + \sqrt{2T f(v_i(t))} \eta_i(t) \quad (7.2)$$

where η_i is a Gaussian white noise.

For the example given in section 5, we can write for (7.1)

$$\dot{u}_i(t) = -E_i(v(t)) - 2T \sinh \frac{2u_i(t)}{a} + 2\sqrt{aT} \cosh \frac{u_i(t)}{a} \eta_i(t) \quad (7.3)$$

A block diagram realization of (7.3) is given in Figure 1.

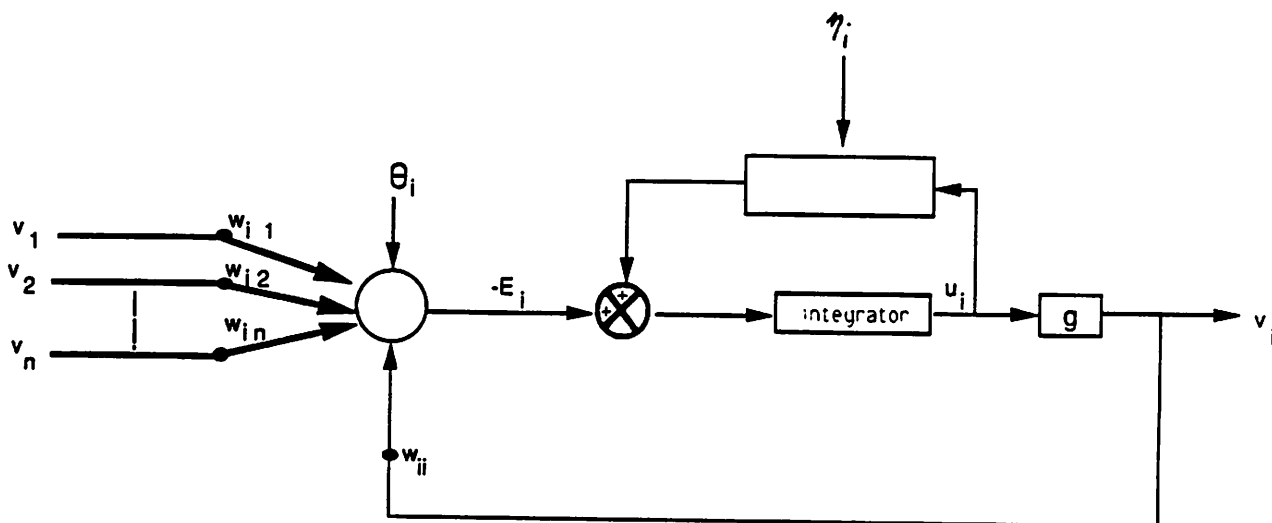


Figure 1: A Node in a Diffusion Machine

8. A Continuous Operating Learning System

Hinton et al. [5] defined a learning problem for Boltzmann machines that can be extended to diffusion machines. Suppose that the n nodes in the network are divided into two groups: *visible* nodes and *hidden* nodes. The space $[0,1]^n$ is correspondingly factored into the cartesian product $V \times H$, where $V = \{v_i : i \text{ visible}\}$ and $H = \{v_i : i \text{ hidden}\}$. Now, suppose that a probability density function \bar{p} is specified in V and we want to find a set of weights w_{ij} so that the stationary distribution of $v_i(t)$ on the visible nodes will be as close to \bar{p} as possible.

We now modify our earlier notation to make the dependencies more explicit. Let $p_0(v, h; w)$ denote the stationary density of the entire network for $v \in V$, $h \in H$, and let w denote the weights. Let the density on V alone be denoted by

$$p_0(v; w) = \int_H p_0(v, h; w) dh \quad (8.1)$$

The problem now is to find w so that $p_0(v; w)$ approximates $\bar{p}(v)$.

Following [5], we use the asymmetric divergence

$$G(w) = \int_V \bar{p}(v) \ln \left[\frac{\bar{p}(v)}{p_0(v; w)} \right] dv \quad (8.2)$$

as a measure of approximation, and shall choose w to minimize $G(w)$. Now,

$$p_0(v, h; w) = \frac{1}{Z(w)} e^{-\frac{1}{T} E(v, h; w)}$$

and

$$Z(w) = \int_{V \times H} e^{-\frac{1}{T} E(v, h; w)} dv dh \quad (8.3)$$

Hence,

$$p_0(v; w) = \frac{\int_H e^{-\frac{1}{T} E(v, h; w)} dh}{\int_{V \times H} e^{-\frac{1}{T} E(v, h; w)} dv dh} \quad (8.4)$$

and

$$\frac{\partial G(w)}{\partial w_{ij}} = \frac{1}{T} \int_{V \times H} \frac{\partial E(v, h; w)}{\partial w_{ij}} \left[p_0(v, h; w) - \bar{p}(v) p_0(h|v; w) \right] dv dh \quad (8.5)$$

where

$$p_0(h|v; w) = \frac{p_0(v, h; w)}{p_0(v; w)} \quad (8.6)$$

If we denote by E_0 the expectation with respect to $p_0(v, h; w)$ and by \bar{E} the expectation with respect to $\bar{p}(v) p_0(h|v; w)$, then we can write

$$\frac{\partial G(w)}{\partial w_{ij}} = \frac{1}{T} \left[E_0 \left[\frac{\partial E}{\partial w_{ij}} \right] - \bar{E} \left[\frac{\partial E}{\partial w_{ij}} \right] \right] \quad (8.7)$$

We first observe that if E is the quadratic function given by (1.3), then

$$\begin{aligned} \frac{\partial E}{\partial w_{ij}} &= -v_i v_j \quad \text{for } i \neq j. \\ &= -\frac{1}{2} v_i^2 \quad \text{for } i = j. \end{aligned}$$

Hence, we can estimate $\frac{\partial G}{\partial w_{ij}}$ by running the network in two modes: (a) *free-running* mode that yields E_0 , and (b) *clamped* mode where for the visible nodes $\{v_i\}$ are found to have a distribution given by \bar{p} to

yield \tilde{E} .

Next, we observe that we can use the Langevin algorithm to minimize $G(w)$. Specifically, we can set

$$d w_{ij}(t) = -G_{ij}(w(t)) dt + \sqrt{2S} d Z_{ij}(t)$$

where $G_{ij} = \frac{\partial G}{\partial w_{ij}}$ comes from the two copies of the network running in two different modes, Z_{ij} are independent Wiener processes, and S is the temperature for the “weight machine” and is different from the temperature T of the networks used to generate G_{ij} . A block diagram of the “learning dynamics” is given in Figure 2.

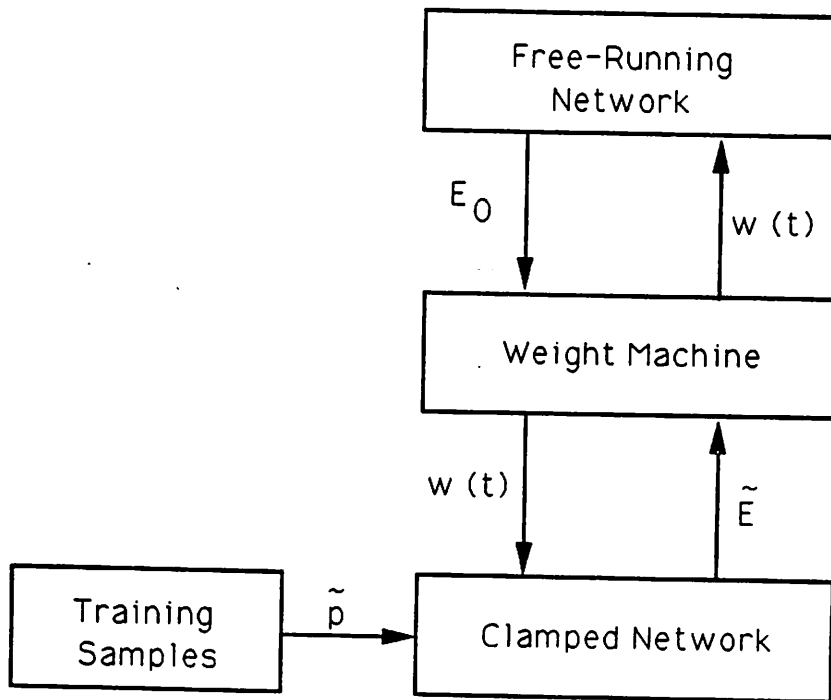


Figure 2: A Continuously Operating Learning System

Intuitively, if we change $W(t)$ slowly in comparison to the dynamics of the two networks, then we should expect the two networks to reach approximate equilibrium before the weights are changed significantly. The continuously operating nature of a diffusion machine in the learning mode makes it a most attractive system.

Conclusion

In this paper we consider a class of stochastic networks which we name *diffusion machines*, that result from a modification of continuous-variable Hopfield networks. We show that by injecting white Gaussian noise in a specific way at each node of a Hopfield network, we obtain a diffusion process with a stationary density of the Boltzmann form. It follows that cooling of the noise sources can be used to achieve annealing, which in turn can be used to obtain a global minimum. As such it is closely related to both the Boltzmann machine and the Langevin algorithm, but superior to both in terms of possible integrated circuit realization [12].

Learning algorithms similar to those proposed from Boltzmann machines are a particularly interesting problem for study. We propose an arrangement consisting of three coupled diffusion machines that perform the functions of training, learning, and weigh-adjustment concurrently in continuous operation. This is in contrast to learning in Boltzmann machines which alternated between a learning phase and an "equilibrating" phase. A substantial speed advantage for diffusion machines is likely.

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